Programmin Languages/Variables and Storage

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Computer Engineering

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Outline

1. Storage
   - Array Variables
2. Semantics of Assignment
3. Variable Lifetime
   - Global Lifetime
   - Local Lifetime
   - Heap Variable Lifetime
   - Dangling Reference and Garbage
   - Persistent Variable Lifetime
4. Commands
   - Assignment
   - Procedure Call
   - Block commands
   - Conditional commands
   - Iterative statements
5. Memory Representation
6. Summary

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- Imperative language variables: variable has a state and value. It can be assigned to different values in same phrase.
- Two basic operations a variable: inspect and update.

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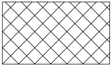

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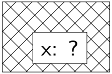


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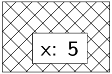


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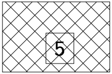
- After the including block terminates, again unallocated

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Total or Selective Update

- Composite variables can be inspected and updated in total or selectively
-

```
struct Complex { double x,y; } a, b;
...
a=b;                    // Total update
a.x=b.y*a.x;            // Selective update
```

- Primitive variables: single cell
  Composite variables: nested cells

Array Variables

Different approaches exist in implementation of array variables:

1. Static arrays
2. Dynamic arrays
3. Flexible arrays

Static arrays

- Array size is fixed at compile time to a constant value or expression.
- C example:

```
  #define MAXELS 100
  int a[10];
  double x[MAXELS*10][20];
}
```

Dynamic arrays

- Array size is defined when variable is allocated. Remains constant afterwards.
- Example: GCC extension (not ANSI!)

```
int f(int n) {
    double a[n]; ...
}
```

- Example: C++ with templates

```
template<class T>  class Array {
      T *content;
  public:
      Array(int s) { content=new T[s]; }
      ~Array()     { delete [] content; }
};
...
Array<int>  a(10);              Array<double> b(n);
```

Flexible arrays

- Array size is completely variable. Arrays may expand or shrink at run time. Script languages like Perl, PHP, Python
- Perl example:

```
@a=(1,3,5);              # array size: 3
print $#a , "\n";        # output: 2 (0..2)
$a[10] = 12;             # array size 11 (intermediate elements ur
$a[20] = 4;              # array size 21
print $#a , "\n";        # output: 20 (0..20)
delete $a[20];           # last element erased, size is 11
print $#a , "\n";        # output: 10 (0..10)
```

- C++ and object orient languages allow overload of [] operator to make flexible arrays possible. STL (Standard Template Library) classes in C++ like vector, map are like such flexible array implementations.

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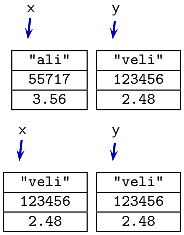
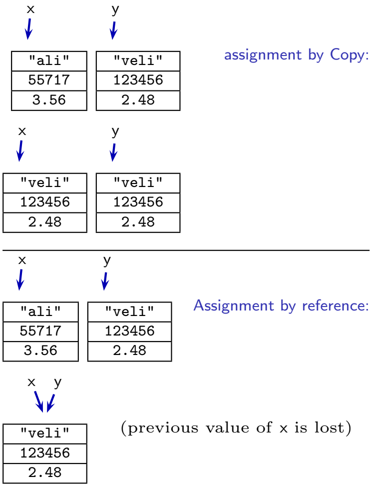


assignment by Copy:

Semantic of assignment in composite variables

- Assignment by Copy vs Reference.
- Copy: All content is copied into the other variables storage. Two copies with same values in memory.
- Reference: Reference of variable is copied to other variable. Two variables share the same storage and values.



- Assignment semantics is defined by the language design
- C structures follows copy semantics. Arrays cannot be assigned. Pointers are used to implement reference semantics. C++ objects are similar.
- Java follows copy semantics for primitive types. All other types (objects) are reference semantics.
- Copy semantics is slower
- Reference semantics cause problems from storage sharing (all operations effect both variables). Deallocation of one makes the other invalid.
- Java provides copy semantic via a member function called `copy()`. Java garbage collector avoids invalid values (in case of deallocation)

Variable Lifetime

- Variable lifetime: The period between allocation of a variable and deallocation of a variable.
- 4 kinds of variable lifetime.
  1. Global lifetime (while program is running)
  2. Local lifetime (while declaring block is active)
  3. Heap lifetime (arbitrary)
  4. Persistent lifetime (continues after program terminates)

Global lifetime

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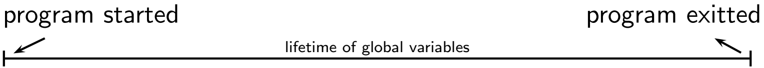

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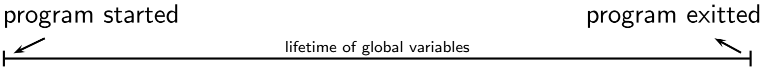

- What are C `static` variables inside functions?

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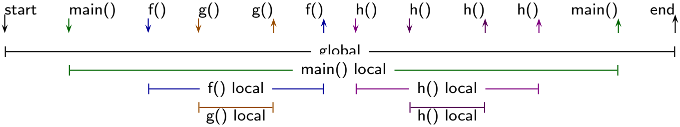
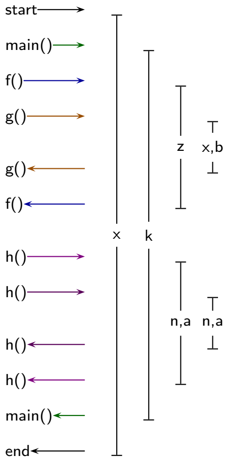
- 

```
double x;
int h(int n) {
   int a;
   if (n<1) return 1
   else return h(n-1);
}
void g() {
   int x;
   int b;
...
}
int f() {
   double z;
   ...
   g();
   ...
}
int main() {
    double k;
    f();
    ...
    h(1);
    ...;
    return 0;
}
```



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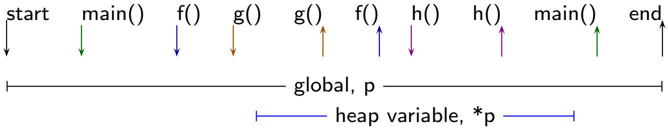
- p and *p are different variables p has pointer type and usually a local or global lifetime, *p is heap variable.
- heap variable lifetime can start or end at anytime.

```
double *p;
int h() { ...
}
void g() { ...
   p=malloc(sizeof(double));
}
int f() { ...
   g(); ...
}
int main() { ...
    f();     ...
    h();     ...;
    free(p); ...
}
```



Dangling Reference

- dangling reference: trying to access a variable whose lifetime is ended and already deallocated.

```
char *p, *q;

p=malloc(10);
q=p;
...
free(q);
printf("%s",p);
```

```
char *f() {
    char a[]="ali";
    ....
    return a;
}
....
char *p;
p=f();
printf("%s",p);
```

- both p's are deallocated or ended lifetime variable, thus dangling reference
- sometimes operating system tolerates dangling references. Sometimes generates run-time erros like "protection fault", "segmentation fault" are generated.

Garbage variables

- garbage variables: The variables with lifetime still continue but there is no way to access.

```
char *p, *q;
...
p=malloc(10);
p=q;
...
```

```
void f() {
    char *p;
    p=malloc(10); ...
    return
}
 ....
f();
```

- When the pointer value is lost or lifetime of the pointer is over, heap variable is unaccessible. (*p in examples)

Garbage collection

- A solution to dangling reference and garbage problem: PL does management of heap variable deallocation automatically. This is called garbage collection. (Java, Lisp, ML, Haskell, most functional languages)
- no call like `free()` or `delete` exists.
- Count of all possible references is kept for each heap variable.
- When reference count gets to 0 garbage collector deallocates the heap variable.
- Garbage collector usually works in a separate thread when CPU is idle.
- Another but too restrictive solution: Reference cannot be assigned to a longer lifetime variable. local variable references cannot be assigned to global reference/pointer.

Persistent variable lifetime

- Variables with lifetime continues after program terminates: file, database, web service object,...
- Stored in secondary storage or external process.
- Only a few experimental language has transparent persistence. Persistence achieved via IO instructions
  C files: `fopen()`, `fseek()`, `fread()`, `fwrite()`
- In object oriented languages; serialization: Converting object into a binary image that can be written on disk or sent to network.
- This way objects snapshot can be taken, saved, restored and object continue from where it remains.

Commands

Expression: program segment with a value. Statement: program segment without a value but with purpose of altering the state. Input, output, variable assignment, iteration...

1. Assignment
2. Procedure call
3. Block commands
4. Conditional commands
5. Iterative commands

Assignment

- C: “`Var = Expr;`”, Pascal “`Var := Expr;`”.
- Evaluates RHS expression and sets the value of the variable at RHS
- `x = x + 1` . LHS `x` is a variable reference (l-value), RHS is the value
- multiple assignment: `x=y=z=0;`
- parallel assignment: (Perl, PHP) `($a,$b) = ($b, $a);`
  `($name, $surname, $no) = ("Onur","Şehitoğlu",55717);`
  Assignment: “reference aggregate” → “value aggregate”
- assignment with operator: `x += 3; x *= 2;`

Procedure call

- Procedure: user defined commands. Pascal: procedure, C: function returning void
- `void` *functname*(*param1*, *param2*, ..., *paramn*)
- Usage is similar to functions but call is in a statement position (on a separate line of program)

Block commands

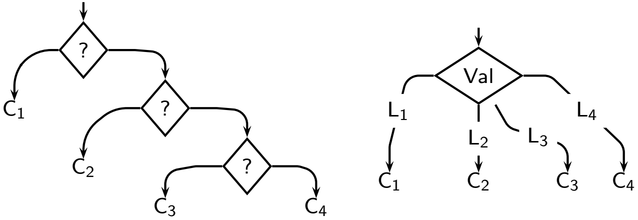
- Composition of a block from multiple statements
- Sequential commands: `{ C`$_1$` ; C`$_2$`; ... ; C`$_n$` }`
  A command is executed, after it finishes the next command is executed,...
- Commands enclosed in a block behaves like single command: "if" blocks, loop bodies,...
- Collateral commands: `{ C`$_1$`, C`$_2$`, ... , C`$_n$` }` (not C ',')!
  Commands can be executed in any order.
- The order of execution is non-deterministic. Compiler or optimizer can choose any order. If commands are independent, effectively deterministic:
  '`y=3 , x=x+1 ;`' vs '`x=3, x=x+1 ;`'
- Can be executed in parallel.

- Concurrent commands: concurrent paradigm languages:
  { C1 | C2 | ... | Cn }
- All commands start concurrently in parallel. Block finishes when the last active command finishes.
- Real parallelism in multi-core/multi-processor machines.
- Transparently handled by only a few languages. Thread libraries required in languages like Java, C, C++.

```
void producer(...) {....}
void collectgarbage(...) {....}
void consumer(...) {....}
int main() {
        ...
        pthread_create(tid1,NULL,producer,NULL);
        pthread_create(tid2,NULL,collectgarbage,NULL);
        pthread_create(tid3,NULL,consumer,NULL);
        ...
}
```

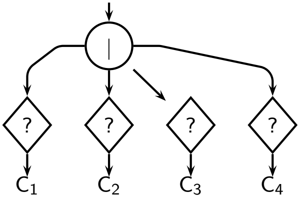
Conditional commands

- Commands to choose between alternative commands based on a condition
- in C : `if (`*cond*`)` C_1 `else` C_2 `;`
  `switch (`*value*`) { case` L_1 `:` C_1 `; case` L_2 `:` C_2 `; ...}`
- `if` commands can be nested for multi-conditioned selection.
- `switch` like commands chooses statements based on a value



- non-deterministic conditionals: conditions are evaluated in collaterally and commands are executed if condition holds.
- hyphotetically:
  `if` ($cond_1$) C_1 `or if` ($cond_2$) C_2 `or if` ($cond_3$) C_3 `;`

  `switch` (val) {
  `case` L_1`:` C_1 | `case` L_2`:` C_2 | `case` L_3`:` C_3 }
- Tests can run concurrently



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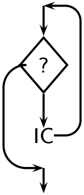

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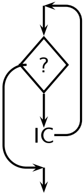
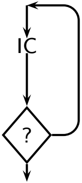


C: `do {...} while (...);`



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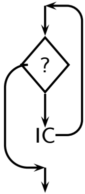

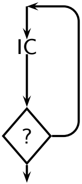

- Another classification: definite vs indefinite iteration

- Definite vs indefinite loops
- Indefinite iteration: Number of iterations of the loop is not known until loop finishes
- C loops are indefinite iteration loops.
- Definite iteration: Number of iterations is fixed when loop started.
- Pascal `for` loop is a definite iteration loop.
  `for i:= k to m do begin .... end;` has $(m - k + 1)$ iterations.
  Pascal forbids update of the loop index variable.
- List and set based iterations: PHP, Perl, Python, Shell

```
$colors=array('yellow','blue','green','red','white');
foreach ($colors as $i) {
    print $i," is a color","\n";
}
```

Memory Representation

- Global variables are kept in fixed region of data segment in memory They are directly accessible
- Heap variables are kept in dynamic region of data segment in memory In a data structure. A memory manager required.
- Local variables are usually kept in run-time stack (Why?)

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- Local variables are usually kept in run-time stack (Why?)
  recursion, each call needs its own set of local variables

Summary

- Variables with storage
- Variable update
- Lifetime: global, local, heap, persistent
- Commands